
Extrapolation of and Code Generation from Communication Traces

Frank Mueller
North Carolina State University



Agenda

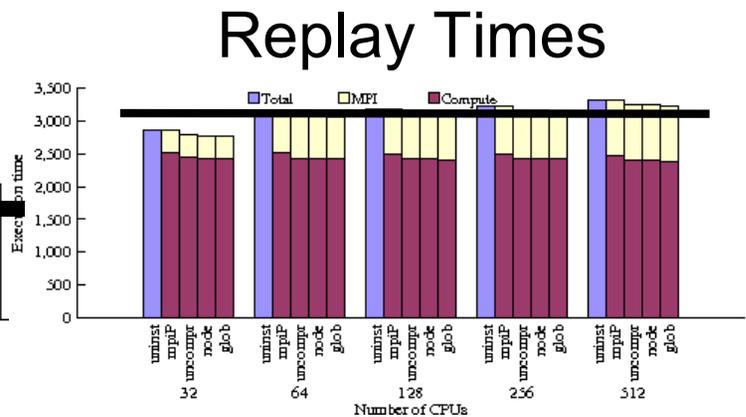
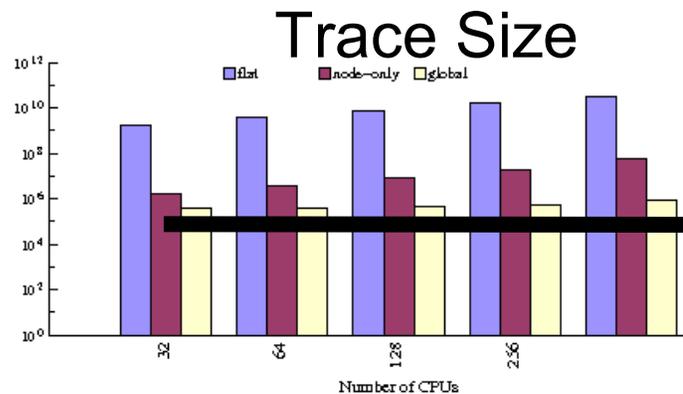
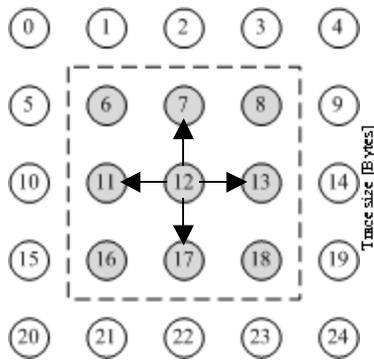
- Overview of ScalaTrace
- Probabilistic Tracing and Replay (ICPP'11)
- Communication Extrapolation of Traces (PPoPP'11)
- Generation of Executable Specifications from Traces (ICS'11)
- Automatic Benchmark Code Generation from Traces

Introduction

- Contemporary HPC Systems
 - Size > 1000 processors
 - take IBM Blue Gene: ~74k nodes, ~300k cores
- Challenges on HPC Systems (large-scale scientific applications)
 - Communication scaling (MPI)
 - Communication analysis
 - Task mapping
- Procurements require performance prediction

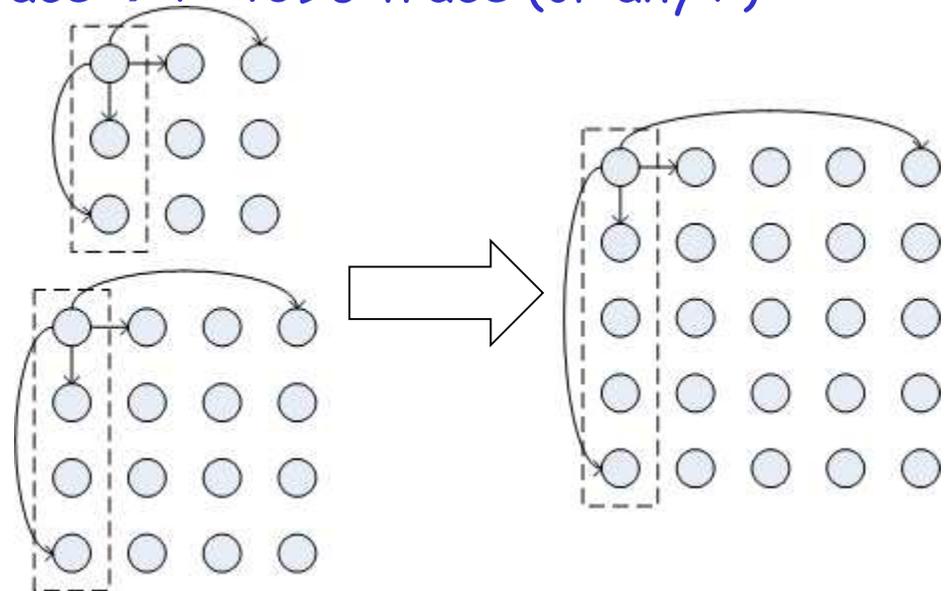
ScalaTrace: Lossless & Scalable Tracing

- Traces communication + I/O of MPI codes via interpositioning
- Trace size: Near constant size, one file represents all nodes
 - Intra-node (loop) & inter-node (task ID) compression of SPMD codes → preserves program structure
 - Location-independent encoding → scales
 - E.g., <10, MPI_Irecv(LEFT), MPI_Isend(RIGHT)>
 - Communication group encoding
 - <dim start_rank iteration_length stride {iteration_length stride}>
- Preserves timing for computation & communication (histograms)



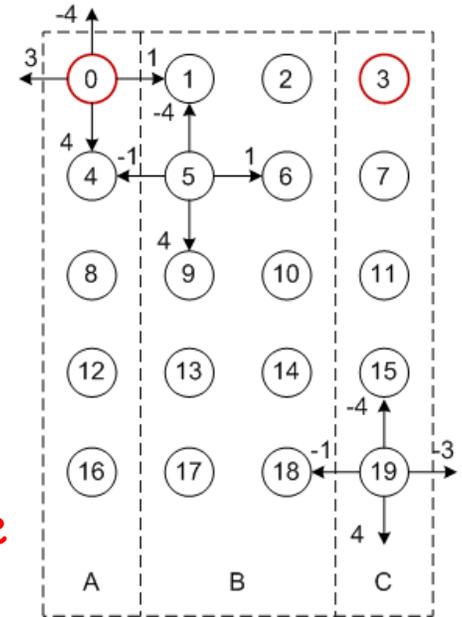
(1) Communication Extrapolation

- Motivation
 - Communication analysis at scale - without running app!
 - Modeling for procurements
 - Extrapolation on a single workstation!
- Idea: synthetically generate communication traces:
 - k small traces from app \rightarrow large traces
 - E.g., $P=8,16,32,64$ nodes trace $\rightarrow P=4096$ trace (or any P)
- Replay large trace/analyze it
- Challenges:
 - Topology detection
 - Message payloads
 - Time extrapolation



Identify & Extrapolate Comm. Topology

- Constrained to row-major stencil/mesh codes
- Algorithm for topology identification:
 - Partition nodes
 - groups according to comm. endpoints
 - Identify critical nodes
 - @ corner/boundary of topo space
 - Calculate boundaries sizes of topological space
 - @ dimension i : $S_i = n_i / n_{i-1}$ (need $i+1$ traces)
- Communication trace extrapolation:
 - Extrapolate comm. params (SRC, DEST, CNT, ...)
 - Extrapolate comm. group



Outline of Algorithm

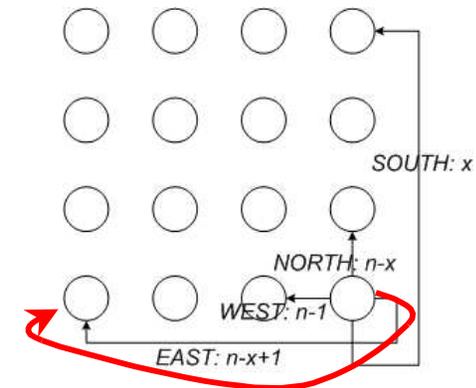
- Input: topology data (x, y, z) from detection phase
- Represent comm. param./group data as $f(x,y,z)$ + unknown coeff.
 - $A^*(xyz) + B^*(xy) + C^*x + D = V$
- Correlate multiple traces & construct set of linear equations
 - $A^*(x_1 y_1 z_1) + B^*(x_1 y_1) + C^*x_1 + D = V_1$
 - $A^*(x_2 y_2 z_2) + B^*(x_2 y_2) + C^*x_2 + D = V_2$
 - $A^*(x_3 y_3 z_3) + B^*(x_3 y_3) + C^*x_3 + D = V_3$
 - $A^*(x_4 y_4 z_4) + B^*(x_4 y_4) + C^*x_4 + D = V_4$

 - Recall: $n_3=xyz, n_2=xy, n_1=x$
- Gaussian Elimination to solve set of equations (solve A B C D)
- With known coefficients, substitute x,y,z at target problem size
- Calculate the desired param. value/comm. group data (V)

Example: Comm. Param. Extrapolation

- Assuming trace inputs for $P = 16, 25, 36, 49$
- Problem: extrapolate relative position of EAST for bottom-right node in a 10×10 mesh

| | |
|-------------------------------------|-----------------------------------|
| $N1=16: x_1=y_1=4, z_1=1, v_1=13$ | $N2=25: x_2=y_2=5, z_2=1, v_2=21$ |
| $N3=36: x_3=y_3=6, z_3=1, v_3=31$ | $N4=49: x_4=y_4=7, z_4=1, v_4=43$ |
| $N_t=100: x_t=y_t=10, z_t=1, v_t=?$ | |



| | | |
|---|----------------------------|---|
| $Ax16 + Bx4x4 + Cx4 + D = 13$ $Ax25 + Bx5x5 + Cx5 + D = 21$ $Ax36 + Bx6x6 + Cx6 + D = 31$ $Ax49 + Bx7x7 + Cx7 + D = 43$ | $A+B = 1$ $C = -1$ $D = 1$ | $v_t = Ax100 + Bx10x10 + Cx10 + D = 91$ |
|---|----------------------------|---|

Example: Comm. Group Extrapolation

- Ranklist of nodes:
 - $\langle \text{dim}, \text{start_rank}, \text{iteration_length}, \text{stride}, \{\text{iteration_length}, \text{stride}\} \rangle$
- Extrapolation performed for
 - start_rank
 - iteration_length
 - stride
- Consider ranklist of Group E

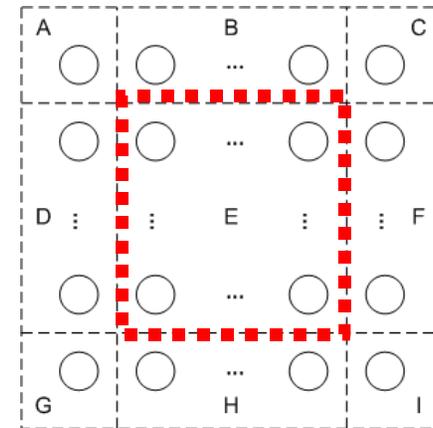
| | |
|---|---|
| $N_1=16: \langle 2 \ 5 \ 2 \ 4 \ 2 \ 1 \rangle$ | $N_2=25: \langle 2 \ 6 \ 3 \ 5 \ 3 \ 1 \rangle$ |
| $N_3=36: \langle 2 \ 7 \ 4 \ 6 \ 4 \ 1 \rangle$ | $N_4=49: \langle 2 \ 8 \ 5 \ 7 \ 5 \ 1 \rangle$ |



| |
|---|
| $\langle 2 \ x+1 \ x-2 \ x \ x-2 \ 1 \rangle$ |
|---|



| |
|--|
| $N_t=100: \langle 2 \ 11 \ 9 \ 10 \ 8 \ 1 \rangle$ |
|--|

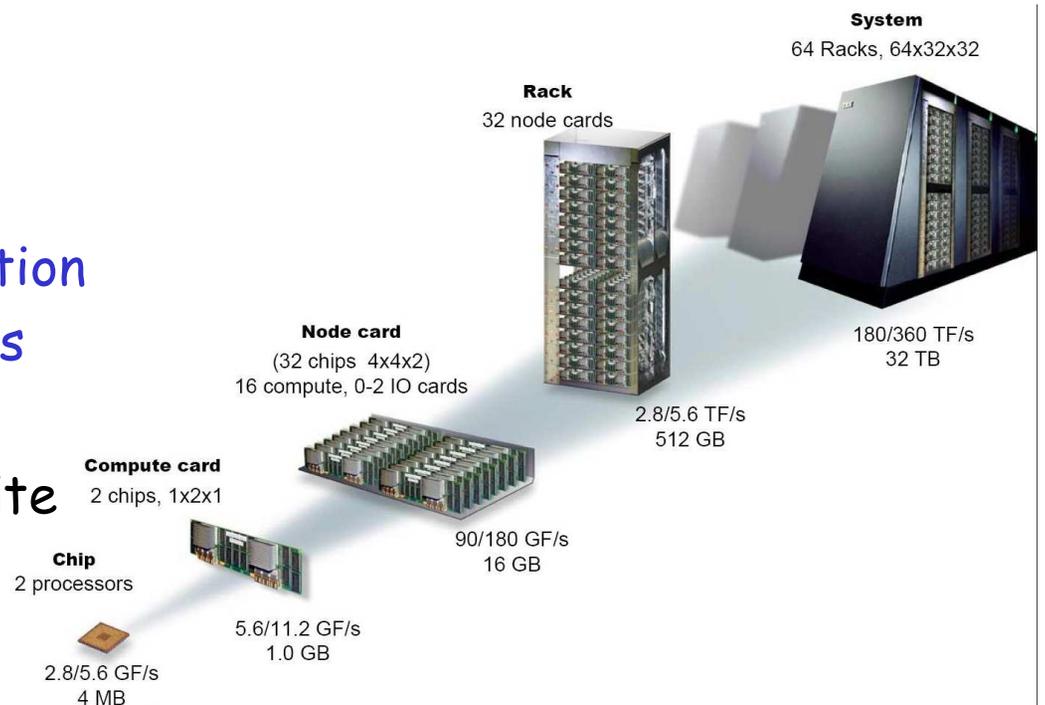


Timing Information Extrapolation

- Match delta times for same computation phase
 - in traces of different size (# nodes)
- Curve fitting to capture variations of delta times
- Currently: 4 types of timing trends
 - **Constant**: $t=f(n)=c \rightarrow$ based on std. dev.
 - **Linear**: $t=f(n)=an+b \rightarrow$ use least-squares method to fit curve
 - **Inverse Proportional**: $t=f(n)=k/n \rightarrow$ std. dev. of $t*n$ (const. k)
 - **Inverse Proportional + Constant**: $t=f(n)=k/n+c$
 \rightarrow least-squares for $t'=tn=cn+k$
- Utilize loosely defined criterion to select best fitting curve
 - **Metric**: std. dev./avg.

Experiment Framework

- Platform: Jugene (IBM BG/P)
 - 73,728 compute nodes and 294,912 cores
 - 2GB memory per node
 - 3D torus and global tree network
- Use subset of nodes
 - base trace generation
 - results verification
- Extrapolation algorithm
 - runs on a single workstation
 - requires only 1-2 seconds
- Experiments with NAS Parallel Benchmark (NPB) suite
 - version 3.3 for MPI

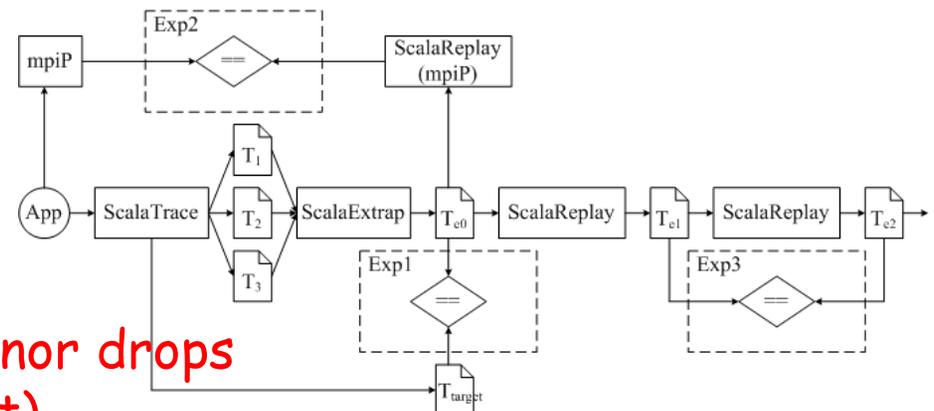


Correctness of Comm. Trace Extrapolation

- for NPB: BT, EP, FT, CG, LU, IS
- Verified in multiple ways:
 1. Extrapolated trace = app trace @ same # nodes size & inputs
 2. Replay extrapolated trace w/ mpiP
 - Aggregate stats = stats from app execution
 3. Replay extrapolated trace w/ ScalaTrace → new trace
 - replay new trace, repeat until fixed point found

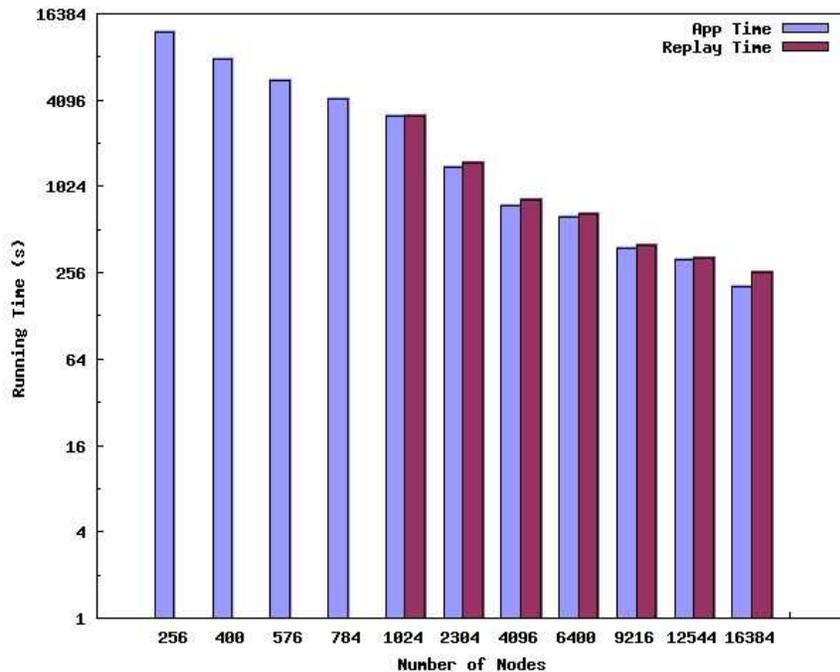
- Results:

- ✓ extrapolated = app traces
- ✓ Aggregate stats match
- ✓ ScalaReplay neither adds nor drops any comm. events (fixpoint)

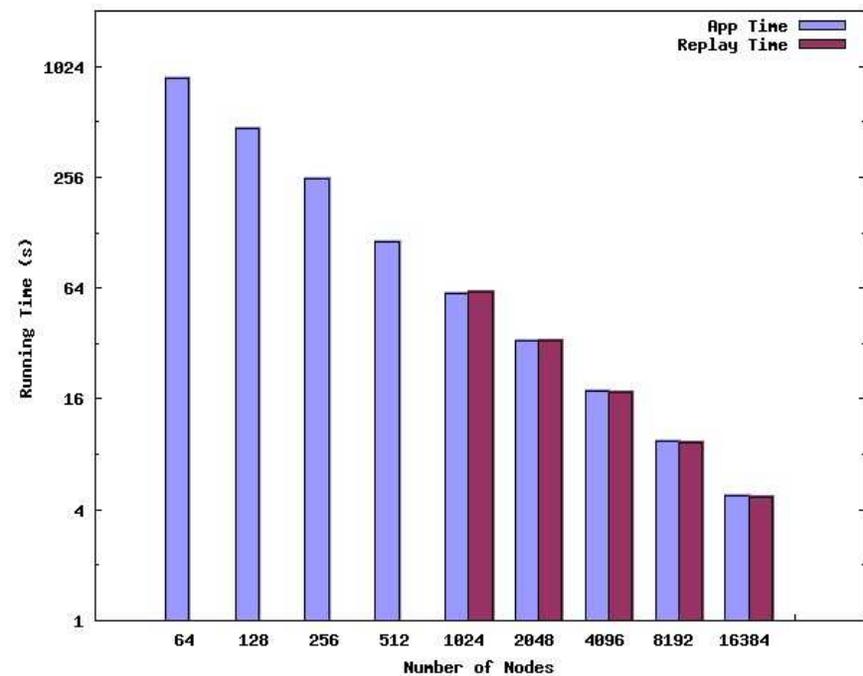


Accuracy of Timing Extrapolation (i)

- for NPB: BT, EP, FT, IS, CG
- up to 16k nodes (not cores)
- $t(\text{extrapolated replay}) = t(\text{app})$
- Accuracy = $|\text{Replay Time} - \text{App Time}| / \text{App Time} \rightarrow \text{generally} > 90\%$



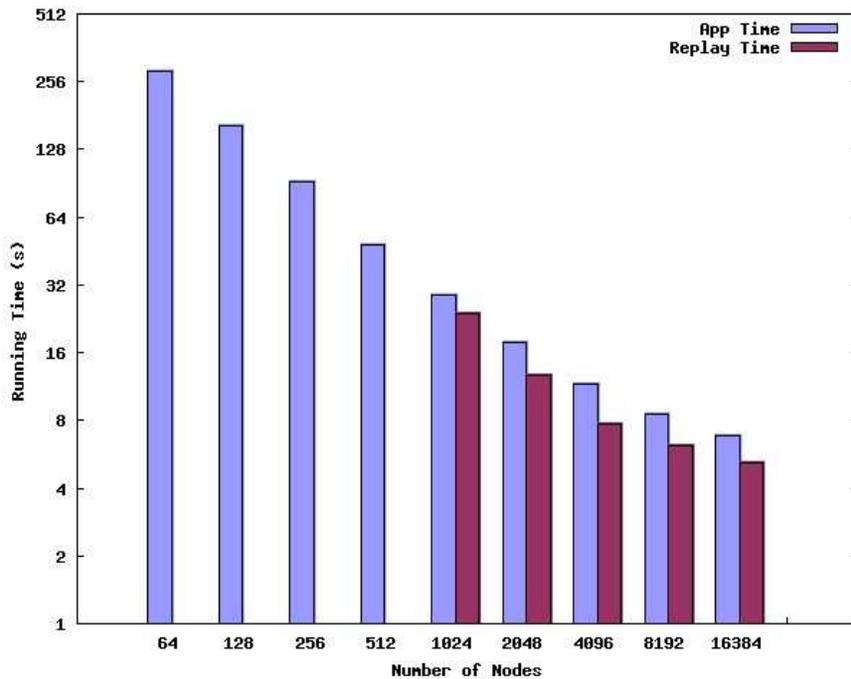
BT (class E)



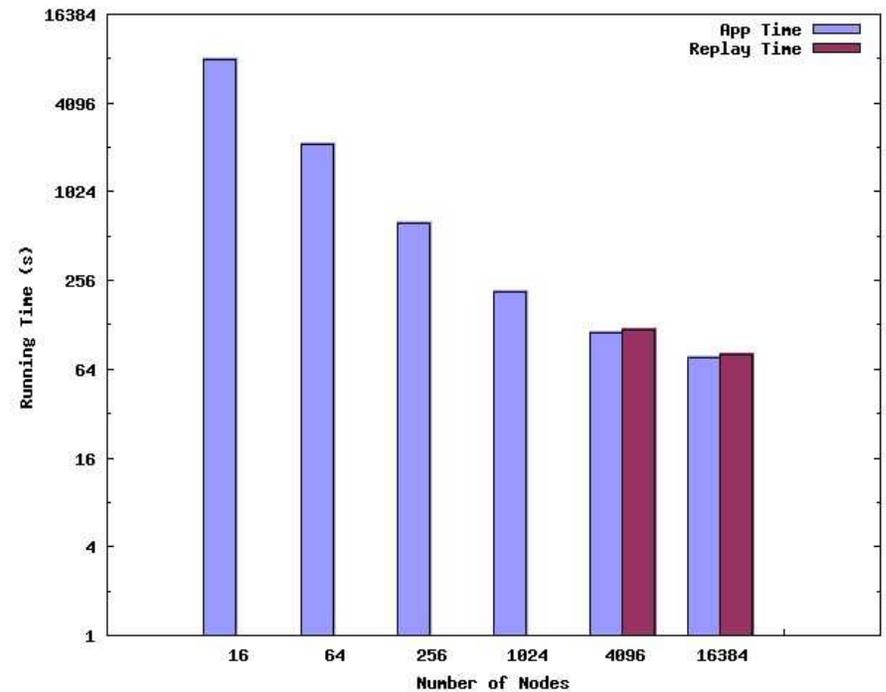
FT (class D)

Accuracy of Timing Extrapolation (ii)

- for NPB: BT, EP, FT, IS, CG
- up to 16k nodes (not cores)
- $t(\text{extrapolated replay}) = t(\text{app})$
- Accuracy = $|\text{Replay Time} - \text{App Time}| / \text{App Time} \rightarrow \text{mostly} > 90\%$



IS (modified input)



CG (class D) square topo

Conclusion for Extrapolation

Contributions:

- Algorithms & techniques for comm. extrapolation, handles
 - trace events
 - execution times
- Based on app runs at smaller scale
- Extrapolation shown to be
 - correct
 - accurate
- Obtains comm. behavior of parallel app at arbitrary scale
 - without actual execution at this scale → unprecedented
- Future Work:
 - Weak scaling

**Extrapolation
not so elusive
anymore**

(2) Communication Benchmark Generation

- *Goal: Generate comm. benchmarks from apps that are*
 - *easy to 1) distribute, 2) use, 3) modify*
- Extracted benchmarks from applications are
 - Performance-accurate
 - Application logic is stripped out
 - Readable, portable, modifiable
 - Collectives are consolidated
 - Nondeterminism has been eliminated
- Target coNCePTuaL: language for rapid generation of network benchmarks
- Compiler + runtime library

```
for(i = 0; i < 10; i++){  
    MPI_Irecv(10, LEFT);  
    MPI_Isend(10, RIGHT);  
    MPI_Waitall();  
}
```



```
For 10 repetitions {  
    All tasks t asynchronously send a  
    10-byte message to task t+1 then  
    all tasks await completion  
}
```

Code Generation from Application Trace

Application → Application Trace → Benchmark in coNCePTual



- All comm. events & computation times generated
- Benchmark contains loop structure
→ easier to read/modify than translation to straight-line code

Consolidating Collectives

- MPI collectives have context sensitive semantic
 - Multiple statements \leftrightarrow single collective
 - Harms readability!

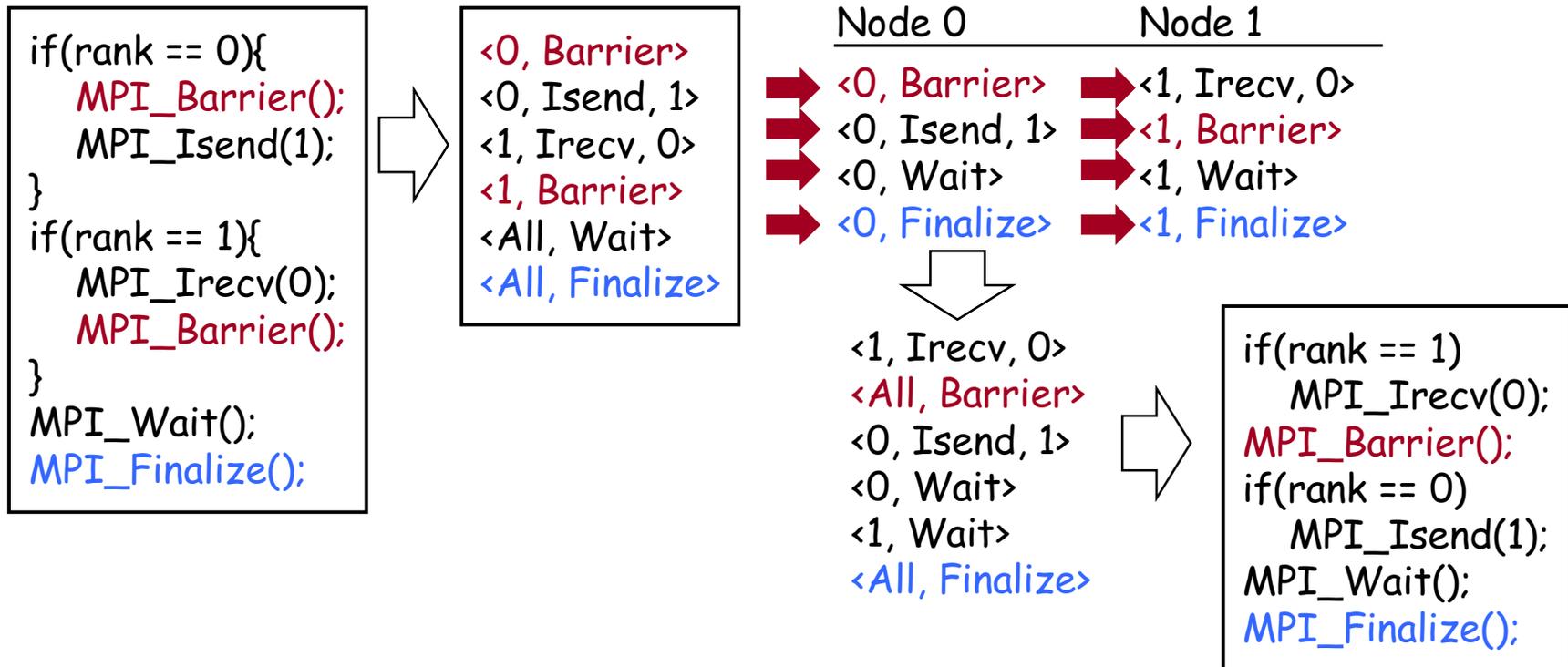
```
if(rank == 1)
  MPI_Barrier();
if(rank == 0){
  MPI_Barrier();
  MPI_Isend(1);
}
if(rank == 1){
  MPI_Irecv(0);
  MPI_Barrier();
}
if(rank == 0)
  MPI_Barrier();
MPI_Wait();
```



```
if(rank == 1)
  MPI_Irecv(0);
  MPI_Barrier();
if(rank == 0)
  MPI_Isend(1);
  MPI_Wait();
```

- Traverse trace with a single node
- Align collectives \rightarrow generate new trace from old one; retain compression
 - Generate RSDs for non-collective events
 - Block & context switch execution at collectives
 - Generate RSD for collective only when last node arrives
 - On-the-fly loop compression

Example: Consolidating Collectives



- On-the-fly loop compression → ensures scalability of trace size

Eliminating Nondeterminism

- Wildcard receives
→ nondeterministic execution

```
MPI_Recv(..., MPI_ANY_SOURCE, ..., status)
if( status.MPI_SOURCE == 0 )
  <Do some LONG-running computation>
else
  <Do some SHORT-running computation>
```

- **Bad performance reproducibility!**
- **Bad readability and modifiability!**
- Replace MPI_ANY_SOURCE with arbitrary valid sender

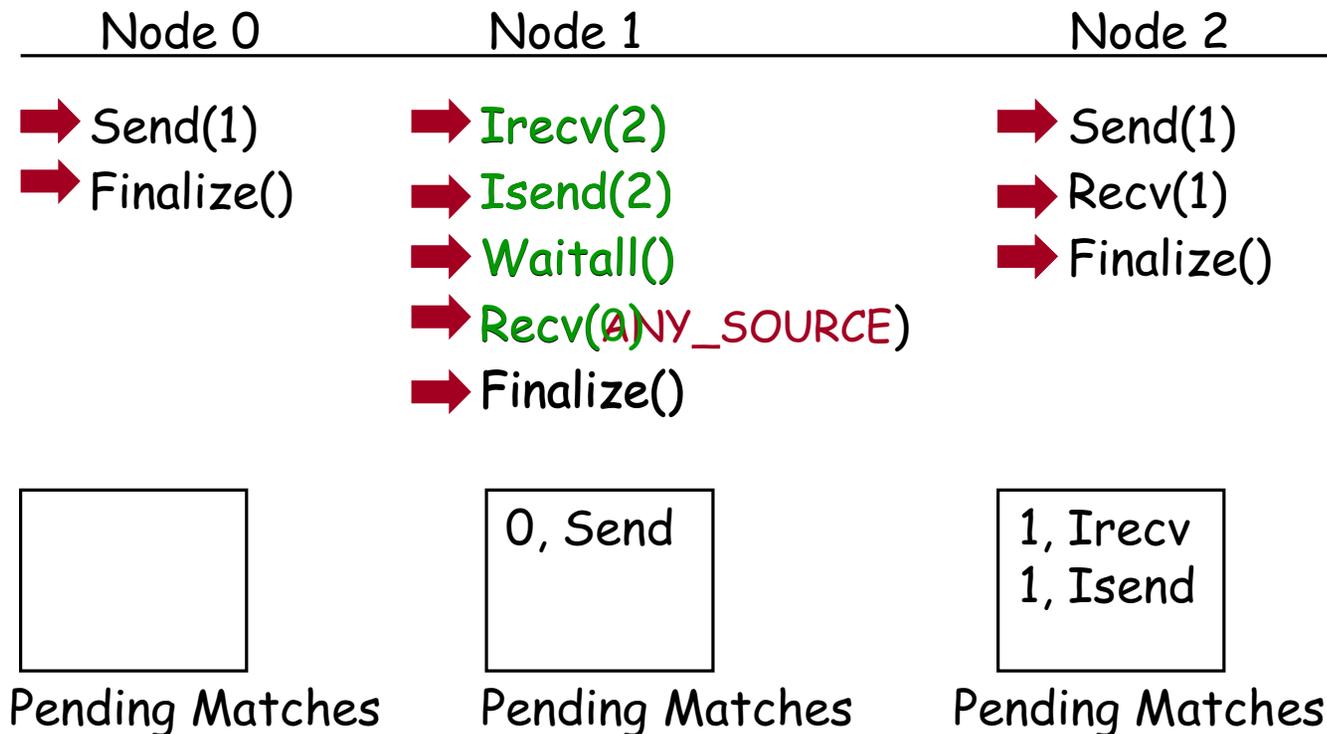
Traverse trace with a single node

Match point-to-point comm.

- Traverse, context sw. if blocked @
 - Blocking Send/Recv
 - Wait, Waitall, etc.
 - Collectives, Finalize
- Switch to node that unblocks current one
- Pair events & unblock an "MPI process" if possible
- Replace MPI_ANY_SOURCE w/ rank of 1st matching sender

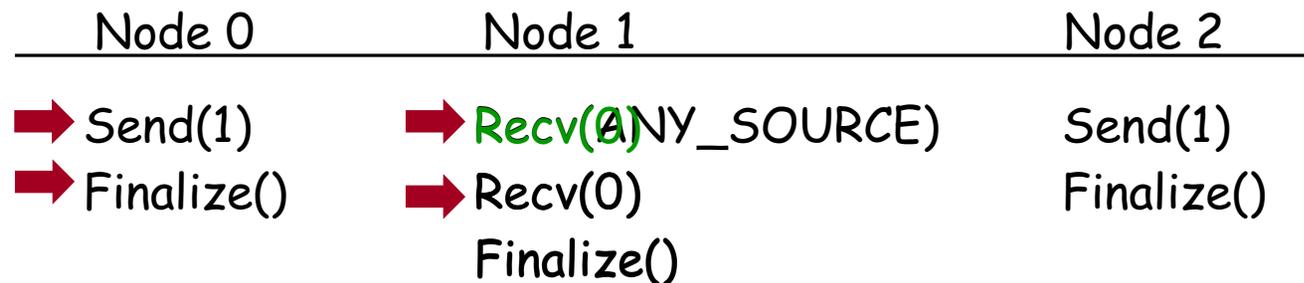
Eliminating Nondeterminism (Example)

- No decompression. No new trace generated.



Eliminating Nondeterminism (Example)

- Deadlock detection: no additional event could be unblocked during last traversal round (across events of all MPI tasks)
- Example: broken program that may deadlock



Cycle Identified!!!

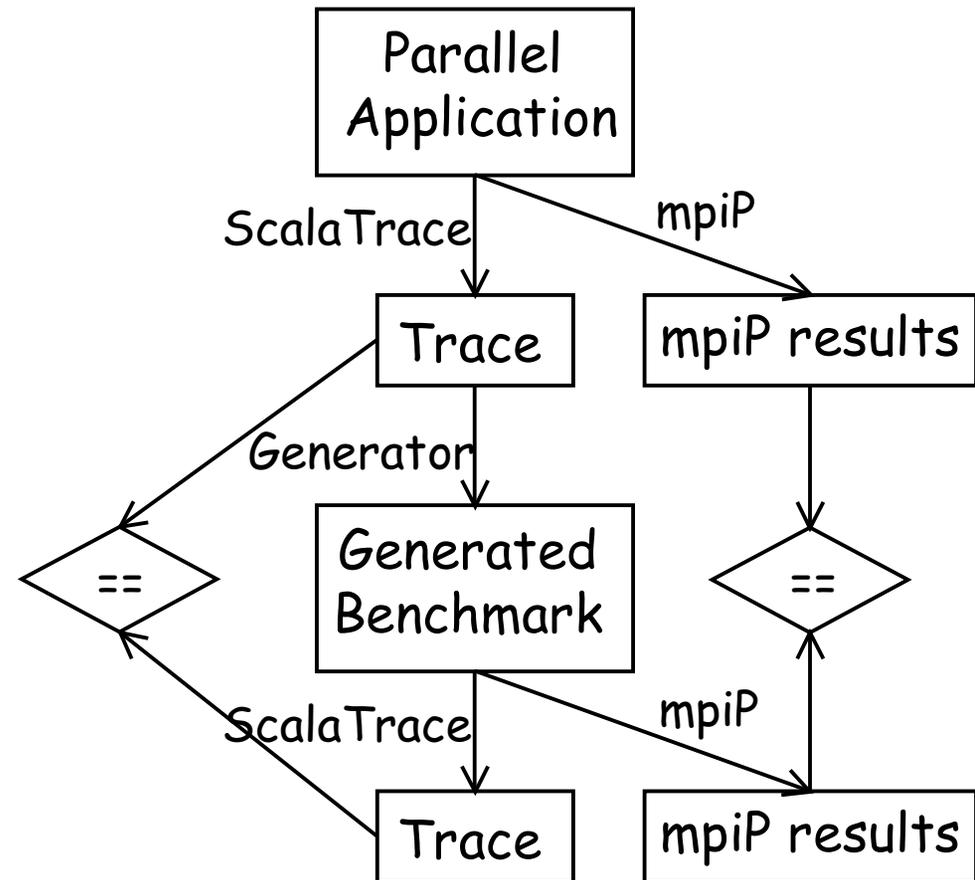
- Above method represents a sufficient condition
 - depends upon the order of traversal
 - If cycle found → deadlock; o/w unknown

Experimental Environment

- Platform:
 - Ocracoke: IBM Blue Gene/L
2,048 nodes, 1GB memory/node
 - ARC: 108 nodes, 1728 cores, 32GB mem/node
 - Single workstation for benchmark generation
- Benchmarks:
 - NAS Parallel Benchmark (NPB V3.3 MPI)
 - Class C inputs, Strong Scaling
 - Sweep3D, Weak Scaling

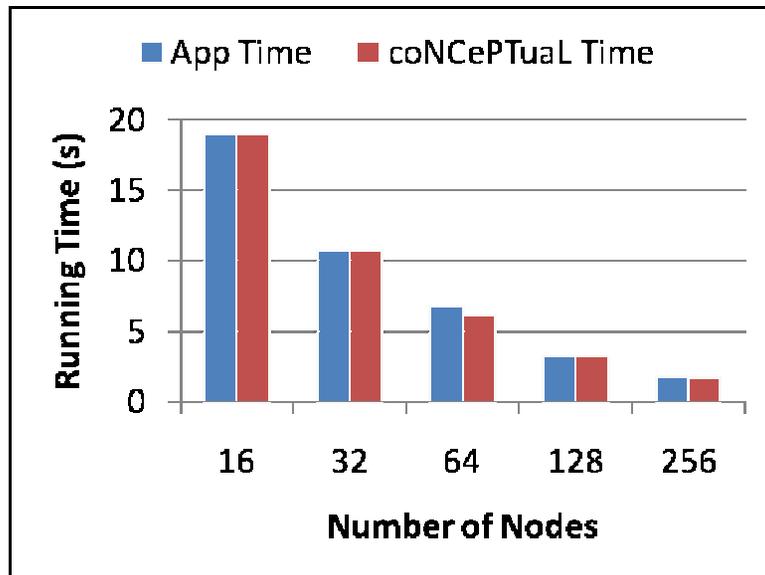
Communication Correctness

- Equivalent traces
 - Communication pattern is preserved
- Equivalent mpiP result
 - Same number of events
 - Same data volume sent
- LU: wildcard receives
- Sweep3D: per-node collective

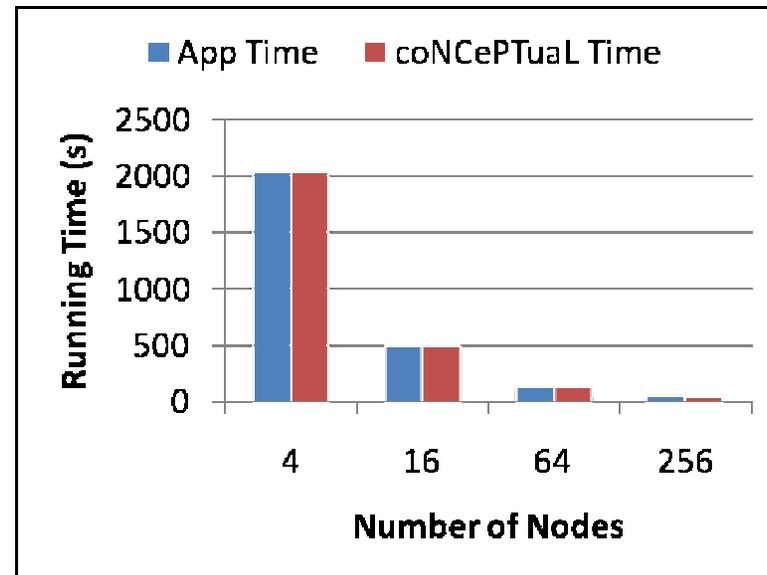


Accuracy of Generated Timings

- Time the application and the generated benchmark, and compare the results
- Mean absolute percentage error is only 2.9% → formula:
$$\frac{|T_{\text{coNCePTuaL}} - T_{\text{app}}|}{T_{\text{app}}} \times 100$$



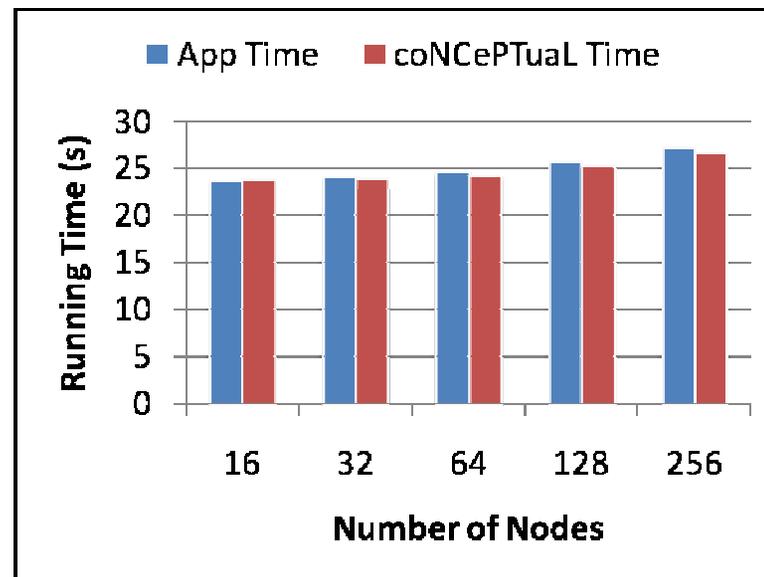
IS



LU

Accuracy of Generated Timings

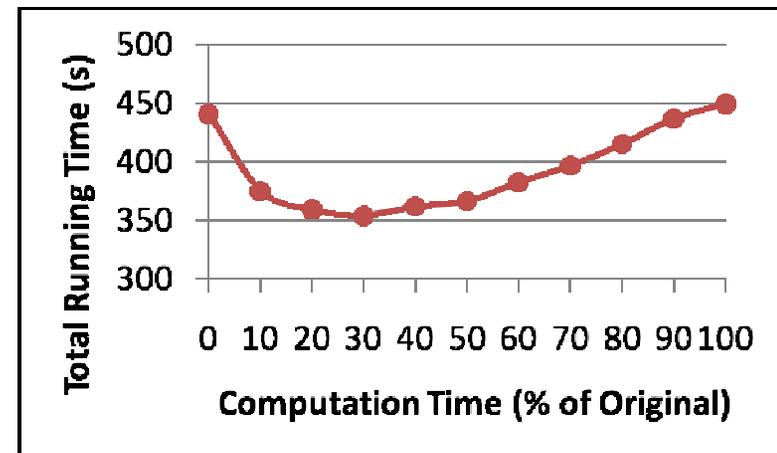
- Time the application and the generated benchmark, and compare the results
- Mean absolute percentage error is only 2.9% → formula:
 $|T_{\text{coNCePTuaL}} - T_{\text{app}}| / T_{\text{app}} \times 100$



Sweep3D: **Weak Scaling**

Applications of the Benchmark Generator

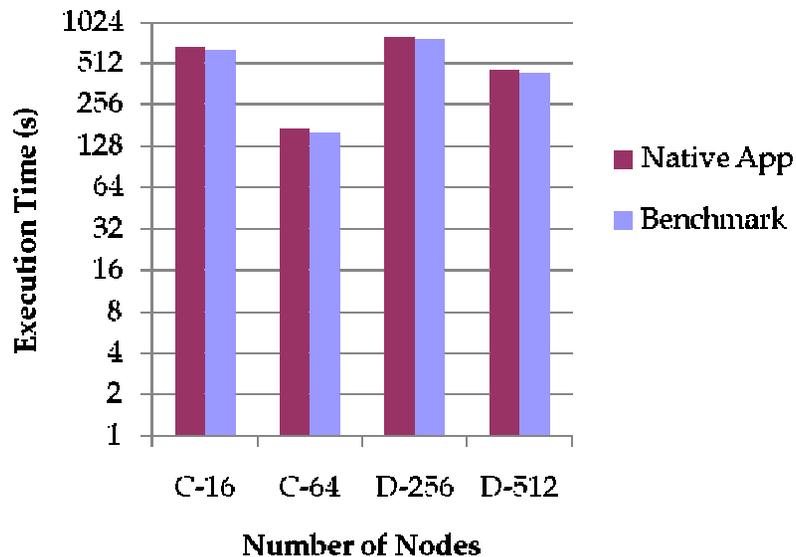
- Determine limits of computation/comm. overlap or effect of computational acceleration (e.g., GPUs)
 - Experiment: ARC cluster, 64 cores, Ethernet, BT b'mark
 - Shorten the spin times gradually
 - 100%: original compute overhead (simulated w/ spin)
 - 0%: no compute overhead → infinitely fast processor
- Best speedup: ~3X
→ overall runtime reduced by 22%
- 0%: network contention or extra memory copies
- Platform-specific result



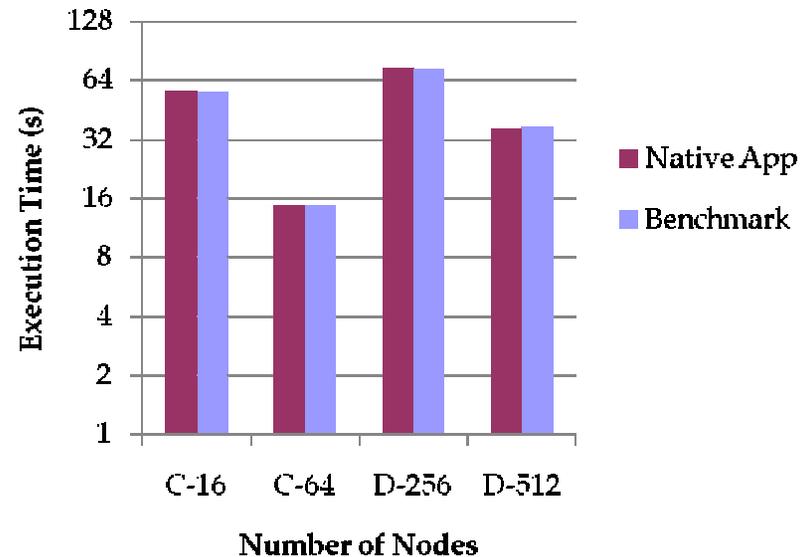
(3) C Benchmark Generation Results – Timing Accuracy

- time from Init → Finalize for app and benchmark
 $|T_{gen} - T_{app}| / T_{app} \times 100$
- Timing accuracy = ~ 6.7% (avg. error)
- ARC cluster (1,7k cores, 16 cores/node Opteron, 32 GB RAM)

LU

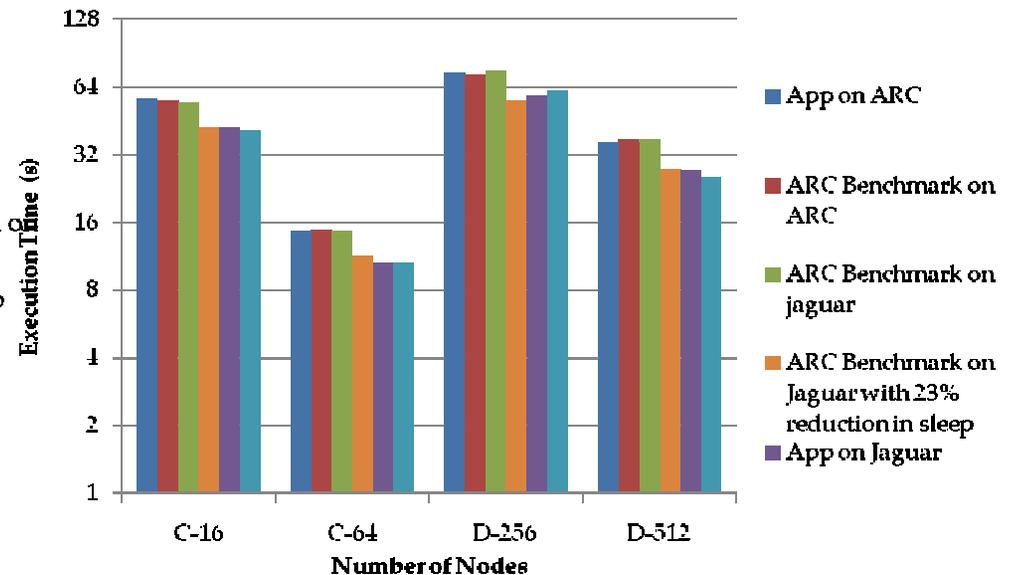
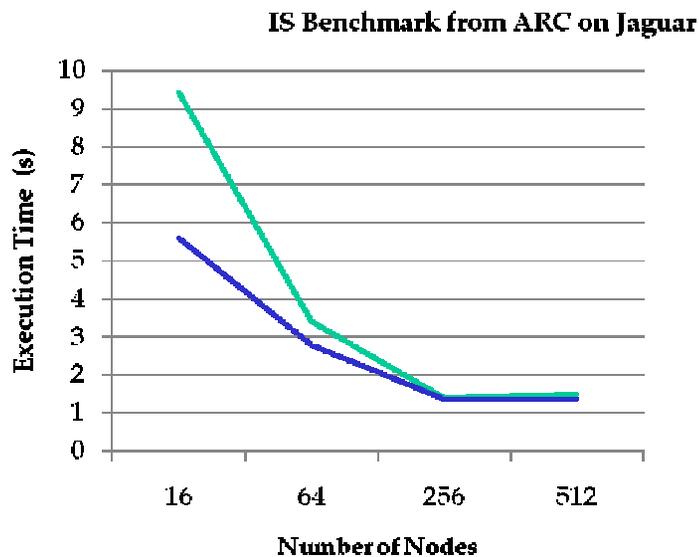


MG



Cross Platform Results: ARC vs. Jaguar

- Jaguar ~23% faster execution on compute kernels
- Resemblance to benchmark obtained on ARC
 - IS: strong scaling reduced per node work
→ close match @ 256 tasks
 - Nearly perfect match after 23% speed correction



Summary

- Trace extrapolation feasible → exascale modeling
- An automatic communication benchmark generation framework
 - ScalaTrace → coNCePTual or C
- Generate benchmarks from real-world apps
 - Ensure performance fidelity, abstract away application logic
 - Readable, portable, modifiable, reproducible
 - Consolidation of collectives
 - Elimination of nondeterminism
 - Obfuscates code → for restricted source code access
 - Facilitate “what-if” analysis

Acknowledgements

- Xing Wu (NCSU, intern @ LANL)
 - Scott Pakin (LANL)
 - Mike Noeth, Prasun Ratn (NCSU, interns @ LLNL)
 - Martin Schulz, Bronis R. de Supinski (LLNL)
 - Karthik Virjayakumar (NCSU, intern @ ORNL)
 - Phil Roth (ORNL)
- 
- 
- 
- best paper [IPDPS'07, JPDC], timed replay [ICS'08], I/O [PDSW'09], extrap [PPoPP'11], gen. specs [ICS'11], prob. Replay [ICPP'11]
 - Code available under BSD license:
moss.csc.ncsu.edu/~mueller/ScalaTrace
 - Funded in part by Humboldt Foundation, NSF 0937908, 0429653, 0410203, CAREER 0237570, 0958311
 - Part of work ... auspices ... U.S. DoE by UC-LLNL under contract No. W-7405-Eng-48 + UT-Batelle, LLC DE-AC05-00OR22725 + Sandia DE-AC52-06NA25396.