



LIBI: The Lightweight Infrastructure-Bootstrapping Infrastructure

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Infrastructure-bootstrapping

Given an infrastructure's binary image(s) and a process/host distribution, start all relevant processes on their respective hosts and propagate necessary startup information.

- Before bootstrapping:
 - Program image on storage device
 - Set of (allocated) computer nodes
- After bootstrapping:
 - Application processes started on computer nodes
 - Application's configuration complete
 - ready for primary operation

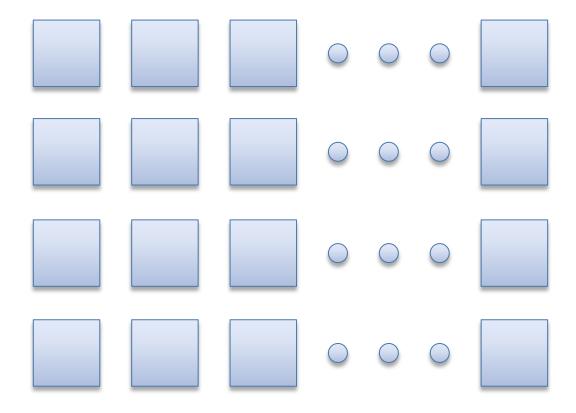


What Infrastructures Need Bootstrapping?

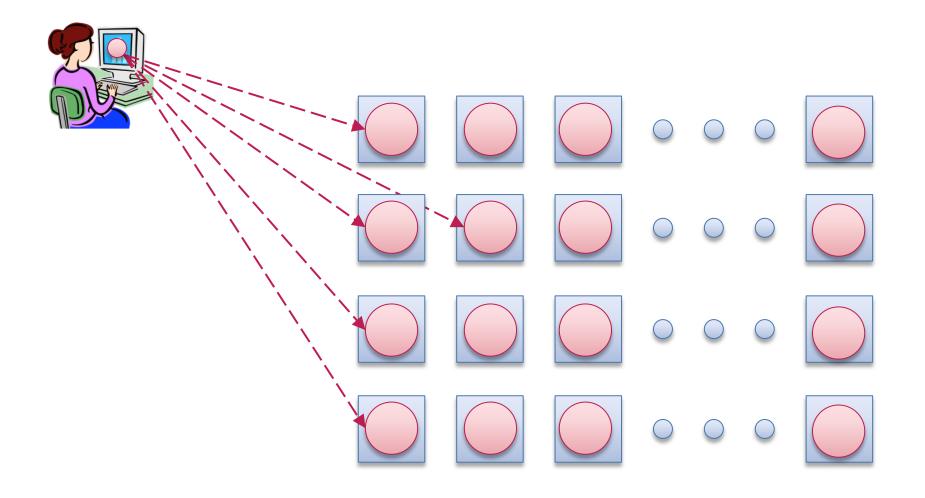
- They all do! Every piece of distributed software needs to be instantiated at some point
 - Applications
 - Tools
 - System services



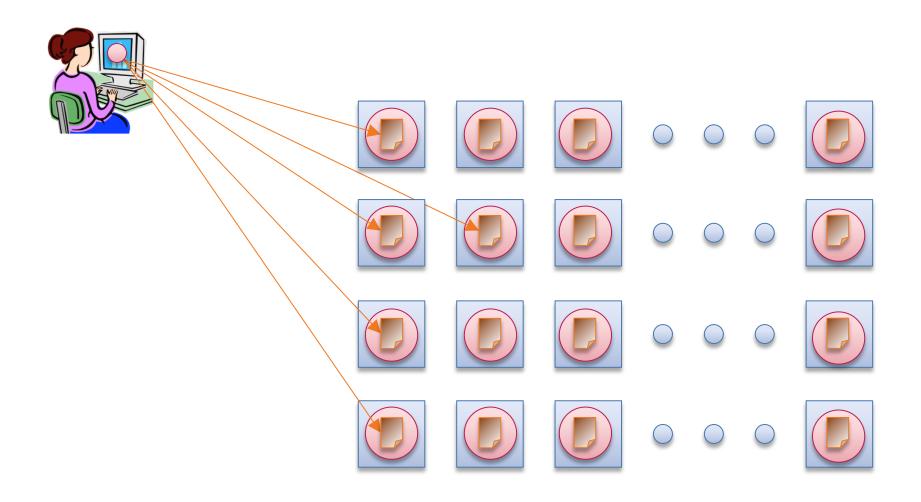




"I gotta get my application up and running!"

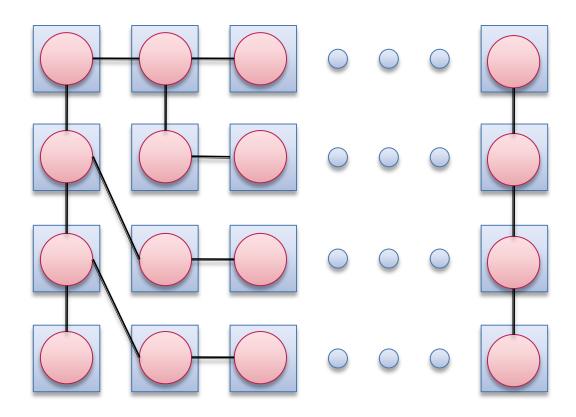


(1) "First, I start all the processes on the appropriate nodes"



(2) "Next, I must disseminate some initialization information"





Bootstrapping is complete when the infrastructure is ready for steady-state usage.

Bootstrapping Operations

Process Launching

▶ Information Dissemination



Bulk Launching Alternatives

- All Strategies Employ Daemons
 - Service daemons or node level agents to start application processes
- Strategies vary along two dimensions
 - Degree of daemon persistence
 - service infrastructure → more persistent
 - application specific → less persistent
 - Daemon interconnection topology
 - simple → less scalabilty
 - hierarchical → more scalability



Degree of (daemon) Persistence

- Persistent daemons, persistent connections
 - MPD
- Persistent daemons, transient connections
 - SLURM, ALPS
- ▶ Transient daemons, transient connections
 - Scela, MRNet default, MRNet on XT



(Daemon) Interconnection Hierarchy

▶ Centralized

▶ Rings

▶ Trees



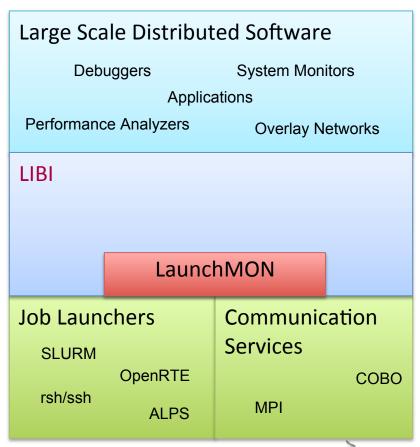
Scalable Bootstrapping Alternatives

- Infrastructure-specific, scalable mechanisms
 - Still limited by sequential operations
 - Not portable to other infrastructures
- Using high-performance resource managers
 - Myriad interfaces
 - No communication facilities
- Generic bootstrapping infrastructures
 - LaunchMON: targets tools with wrapper for existing RMs



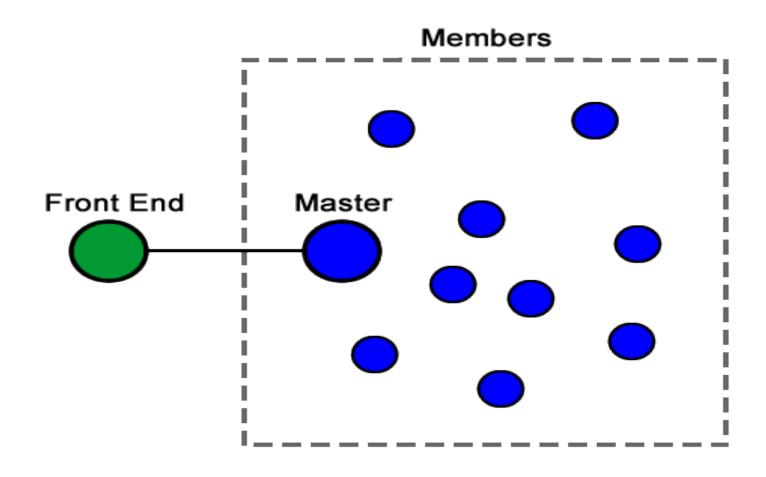
LIBI Approach

- LIBI: Lightweight infrastructurebootstrapping infrastructure
 - Generic service for scalable distributed software infrastructure bootstrapping
 - Process launch
 - Scalable, low-level collectives





LIBI Architecture



LIBI API

- session: set of processes (to be) deployed
 - session master manages other members
 - session front-end interacts with session master
 - LIBI currently supports only master/member communication
- host-distribution: where to create processes
 - <hostname, num-processes>
- process distribution: how/where to create processes
 - <session-id, executable, arguments, host-distribution, environment>

LIBI API (cont'd)

launch (process-distribution-array)

instantiate processes according to input distributions

[send|recv]UsrData(session-id, msg)

communicate between front-end and session master

broadcast(), scatter(), gather(), barrier()

communicate amongst session members



Example LIBI Front-end

```
front-end(){
  LIBI fe init();
  LIBI fe createSession(sess);
  proc dist req t pd;
  pd.sessionHandle = sess;
  pd.proc path = get ExePath();
  pd.proc argv = get ProgArgs();
  pd.hd = get HostDistribution();
  LIBI fe launch(pd);
  //test broadcast and barrier
  LIBI fe sendUsrData(sess1, msg, len );
  LIBI fe recvUsrData(sess1, msg, len);
  //test scatter and gather
  LIBI fe sendUsrData(sess1, msg, len );
  LIBI fe recvUsrData(sess1, msg, len);
  return 0;
```

Example LIBI-launched Application

```
session_member() {
   LIBI_init();

   //test broadcast and barrier
   LIBI_recvUsrData(msg, msg_length);
   LIBI_broadcast(msg, msg_length);
   LIBI_barrier();
   LIBI_sendUsrData(msg, msg_length)

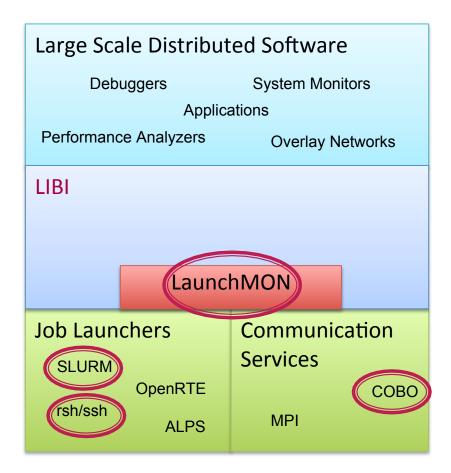
   //test scatter and gather
   LIBI_recvUsrData(msg, msg_length);
   LIBI_scatter(msg, sizeof(rcvmsg), rcvmsg);
   LIBI_gather(sndmsg, sizeof(sndmsg), msg);
   LIBI_sendUsrData(msg, msg_length);

LIBI_finalize();
}
```



LIBI Implementation Status

- LaunchMON-based runtime
 - Tested SLURM or rsh launching
 - COBO PMGR service
- Rsh-based default
 - Pluggable launch topologies
 - Devised a provably optimal algorithm!



Optimal Launching Topology

Assumptions

- Homogenous computing environment
 - All nodes have the same computational power
 - Constant wait time between each local launch command
 - Constant remote launch time
 - physical network topology?
 - file system (and other resource) contention?

Algorithm Overview

- Inspired by Park et al's optimal multicast tree [ICPP '96]
- Pick first node as root
- For every subsequent node, place at minimal launch point



Algorithm for Optimal Launch Topology

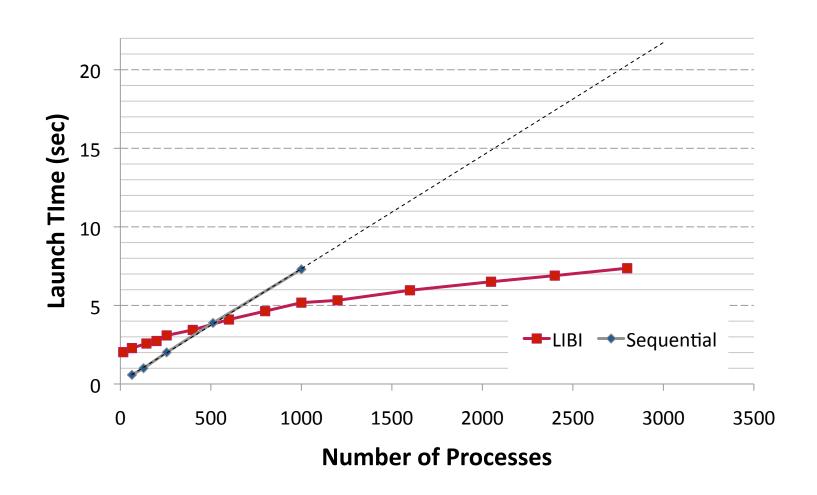
```
find optimal topology ( node list, model params ) {
  dequeue list head, set as root of tree
  compute root's "ready time"
  while( node list not empty ) {
     dequeue list head
      add node to tree at smallest "ready time" node
      compute node's "ready time"
      recompute parent's "ready time"
```

Performance Results

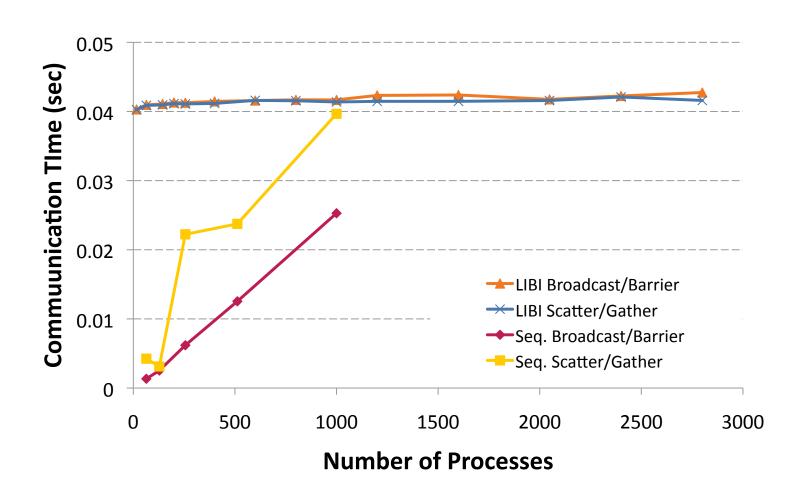
- Focus on task launching
 - Lots of data available for communication topologies
- MRNet Start-up improvements



LIBI Microbenchmark Results



LIBI Microbenchmark Results



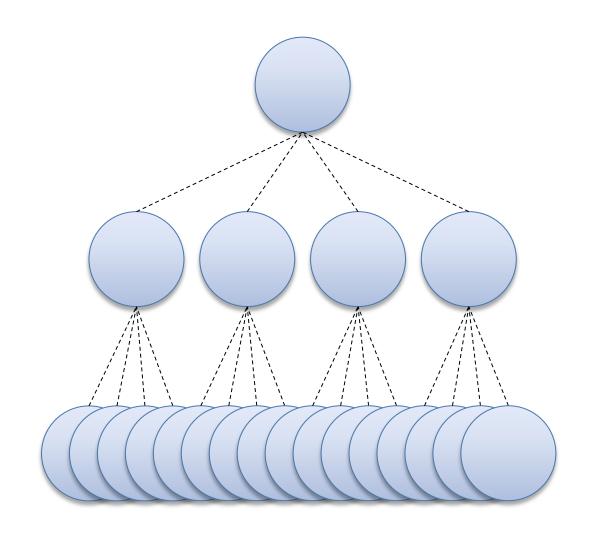
MRNet/LIBI Integration

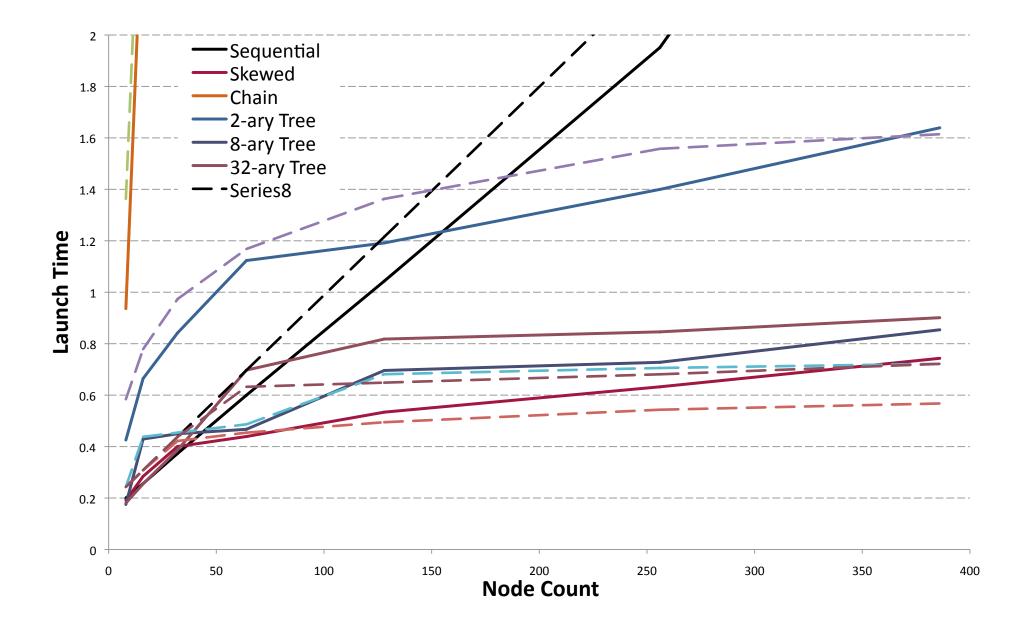
▶ MRNet uses LIBI to launch all MRNet processes

- Parse topology file and setup/call LIBI launch ()
- Session front-end gathers/scatters startup information
 - Parent listening socket (IP/port)

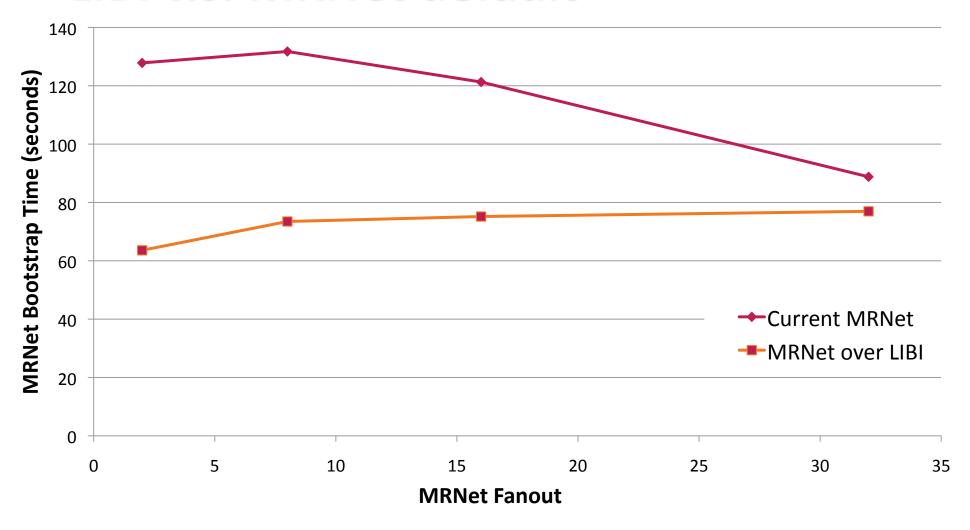
MRNet Sequential-ish Bootstrapping

- Parent creates children
 - ▶ Local → fork()/exec()
 - ▶ Remote → rsh-based mechanism
- Integrated instantiation and information propagation
- MRNet's "standard"

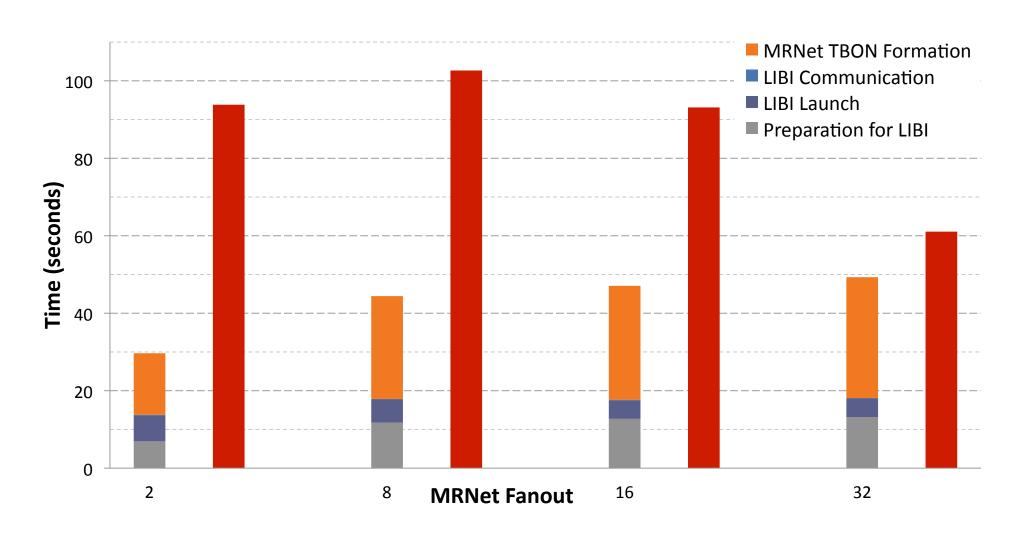




LIBI v.s. MRNet default



LIBI v.s. MRNet default (broken down)



Future Research and Development

- Optimal Launch Topology
 - Performance analysis
 - Scalability bounds
 - Optimization heuristics
 - Impact of resource contention and other simplifications
- Mechanisms to alleviate file system contention
 - Like our scalable binary relocation service
- More flexible process and host distributions
 - Instantiating different images in same session
 - Integrating allocating and launching
- Integrated scalable communication infrastructure
- Refactoring LaunchMON